

# Advanced Data Modeling

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- ◆ First-order logic. Syntax and semantics.
- ◆ Herbrand interpretations;
- ◆ Clauses and goals;
- ◆ Datalog.

First-order signature  $\Sigma$  consists of

- ◆ *con* — the set of constants of  $\Sigma$ ;
- ◆ *fun* — the set of function symbols of  $\Sigma$ ;
- ◆ *rel* — the set of relation symbols of  $\Sigma$ .

Term of  $\Sigma$  with variables in  $X$ :

1. Constant  $c \in con$ ;
2. Variable  $v \in X$ ;
3. If  $f \in fun$  is a function symbol of arity  $n$  and  $t_1, \dots, t_n$  are terms, then  $f(t_1, \dots, t_n)$  is a term.

- ◆ A term is ground if it has no variables
- ◆  $var(t)$  — the set of variables of  $t$

Abstract notation:

- ◆  $a, b, c, d, e$  for constants;
- ◆  $x, y, z, u, v, w$  for variables;
- ◆  $f, g, h$  for function symbols;
- ◆  $p, q$  for relation symbols,

Example:  $f(x, g(y))$ .

Concrete notation: teletype font for everything.

Variable names start with upper-case letters.

Example: `likes(john, Anybody)`.

- ◆ Atomic formulas, or atoms  $p(t_1, \dots, t_n)$ .
- ◆  $(A_1 \wedge \dots \wedge A_n)$  and  $(A_1 \vee \dots \vee A_n)$
- ◆  $(A \rightarrow B)$  and  $(A \leftrightarrow B)$
- ◆  $\neg A$
- ◆  $\forall v A$  and  $\exists v A$

- ◆ Substitution  $\theta$  : is any mapping from the set  $V$  of variables to the set of terms such that there is only a finite number of variables  $v \in V$  with  $\theta(v) \neq v$ .
- ◆ Domain  $dom(\theta)$ , range  $ran(\theta)$  and variable range  $vran(\theta)$ :
  - $dom(\theta) = \{v \mid v \neq \theta(v)\}$ ,
  - $ran(\theta) = \{t \mid \exists v \in dom(\theta)(\theta(v) = t)\}$ ,
  - $vran(\theta) = var(ran(\theta))$ .
- ◆ Notation:  $\{x_1 \mapsto t_1, \dots, x_n \mapsto t_n\}$
- ◆ empty substitution  $\{\}$

Application of a substitution  $\theta$  to a term  $t$ :

- ◆  $x\theta = \theta(x)$
- ◆  $c\theta = c$
- ◆  $f(t_1, \dots, t_n)\theta = f(t_1\theta, \dots, t_n\theta)$



A Herbrand interpretation of a signature  $\Sigma$  is any set of ground atoms of this signature.

1. If  $A$  is atomic, then  $I \models A$  if  $A \in I$
2.  $I \models B_1 \wedge \dots \wedge B_n$  if  $I \models B_i$  for all  $i$
3.  $I \models B_1 \vee \dots \vee B_n$  if  $I \models B_i$  for some  $i$
4.  $I \models B_1 \rightarrow B_2$  if either  $I \models B_2$  or  $I \not\models B_1$
5.  $I \models \neg B$  if  $I \not\models B$
6.  $I \models \forall xB$  if  $I \models B\{x \mapsto t\}$  for all ground terms  $t$  of the signature  $\Sigma$
7.  $I \models \exists xB$  if  $I \models B\{x \mapsto t\}$  for some ground term  $t$  of the signature  $\Sigma$

- ◆ Literal is either an atom or the negation  $\neg A$  of an atom  $A$ .
- ◆ Positive literal: atom
- ◆ Negative literal: negation of an atom
- ◆ Complimentary literals:  $A$  and  $\neg A$
- ◆ Notation:  $L$

Clause: (or normal clause) formula  $L_1 \wedge \dots \wedge L_n \rightarrow A$ ,

where

- ◆  $n \geq 0$ , each  $L_i$  is a literal and  $A$  is an atom.
- ◆ Notation:  $A :- L_1 \wedge \dots \wedge L_n$  or  $A :- L_1, \dots, L_n$
- ◆ Head: the atom  $A$ .
- ◆ Body: The conjunction  $L_1 \wedge \dots \wedge L_n$
- ◆ Definite clause: all  $L_i$  are positive
- ◆ Fact: clause with empty body

Clause	Class
<code>lives(Person, sweden) :- sells(Person, wine, Shop), not open(Shop,saturday)</code>	normal
<code>spy(Person) :- russian(Person)</code>	definite
<code>spy(bond)</code>	fact

- ◆ Goal (also normal goal) is any conjunction of literals  
 $L_1 \wedge \dots \wedge L_n$
- ◆ Definite goal: all  $L_i$  are positive
- ◆ Empty goal  $\square$ : when  $n = 0$