Relational Data Model
Overview

- Relational data model;
- Tuples and relations;
- Schemas and instances;
- Named vs. unnamed perspective;
- Relational algebra;
<table>
<thead>
<tr>
<th>Player</th>
<th>Birth Year</th>
</tr>
</thead>
<tbody>
<tr>
<td>Andy</td>
<td>1980</td>
</tr>
<tr>
<td>Wim</td>
<td>1975</td>
</tr>
<tr>
<td>Liam</td>
<td>1985</td>
</tr>
<tr>
<td>Mike</td>
<td>1988</td>
</tr>
<tr>
<td>Bert</td>
<td>1971</td>
</tr>
</tbody>
</table>
Observations

- The **rows** of the table contain pairs occurring in the relation **player**.
- There are two **columns**, labeled respectively by “name” and “birth year”.
- The **values** in each column belong to different **domains** of possible values.
How to specify a relation

1. specifying the names of the columns (also called fields or attributes);

2. specifying a domain of possible values for each column;

3. enumerate all tuples in the relation.

(1)–(2) refer to the schema of this relation, while (3) to an instance.
Domains and attributes

- A set of domains $\mathcal{D}_1, \mathcal{D}_2$ (sets of values);

- A set of corresponding domain names $d_1, d_2, \ldots$

- A set of attributes $a_1, a_2, \ldots$
Tuple: any finite sequence \((v_1, \ldots, v_n)\). \(n\) is the **arity** of this tuple.

Relation schema:

\[ r(a_1:d_1, \ldots, a_n:d_n) \]

where \(n \geq 0\), \(r\) is a relation name, \(a_1, \ldots, a_n\) are distinct attributes, \(d_1, \ldots, d_n\) are domain names.

Relation instance:

finite set of tuples \((v_1, \ldots, v_n)\) of arity \(n\) such that \(v_i \in D_i\) for all \(i\).
Observation

1. The attributes in each column must be unique.

2. A relation is a set. Therefore, when we represent a relation by a table, the order of rows in the table does not matter.

Let us add to this:
3. The order of attributes does not matter.
New notation for tuples

- A tuple is a set of pairs \( \{(a_1, v_1), \ldots, (a_n, v_n)\} \) denoted by \( \{a_1=v_1, \ldots, a_n=v_n\} \).

- Let \( d_1, \ldots, d_n \) be domain names and \( \mathcal{D}_1, \ldots, \mathcal{D}_n \) be the corresponding domains.

- The tuple \textbf{conforms} to a relation schema \( r(a_1:d_1, \ldots, a_n:d_n) \) if \( v_i \in D_i \) for all \( i \).
Relational data are structured

Note that in the relational data model tuples stored in a table are **structured**:

- all tuples conform to the same relation schema;
- the values in the same column belong to the same domain.

**Untyped perspective:** there is a single domain, so the second condition can be dropped.
**Typed or untyped?**

- Consider the relation admire:

<table>
<thead>
<tr>
<th>admirer</th>
<th>admired</th>
</tr>
</thead>
<tbody>
<tr>
<td>wim</td>
<td>andy</td>
</tr>
<tr>
<td>mike</td>
<td>wim</td>
</tr>
<tr>
<td>liam</td>
<td>andy</td>
</tr>
<tr>
<td>liam</td>
<td>arsenal</td>
</tr>
</tbody>
</table>
Database schema and instance


- Relational database instance conforming to a relational database schema $S$:
  - a mapping $I$ from the relation names of $S$ to relation instances such that
    - for every relation schema $r(a_1:d_1, \ldots, a_n:d_n)$ in $S$ the relation instance $I(r)$ conforms to this relation schema.
No attributes

- a tuple is simply a sequence \((v_1, \ldots, v_n)\) of values.

- The components of tuples can therefore be identified by their position in the tuple.
From Unnamed to Named Perspective

- Introduce a collection of attributes #1,#2, …,

- identify tuple \((v_1, \ldots, v_n)\) with the tuple 
  \{ #1 = v_1, \ldots, #n = v_n \}.

- Likewise, identify relation schema \(r(d_1, \ldots, d_n)\) with 
  \(r(\#1:d_1, \ldots, \#n:d_n)\).
1. Can **define** new relations from existing ones;
2. Uses a collection of **operations** on relations to do so.
\{ (v_{11}, \ldots, v_{1n}), \\
\ldots \\
(\nu_{k1}, \ldots, \nu_{kn}) \}
WeST

\[ R_1 \cup R_2 = \{ (c_1, \ldots, c_k) \mid (c_1, \ldots, c_k) \in R_1 \text{ or } (c_1, \ldots, c_k) \in R_2 \} \]
Set difference

\[ R_1 - R_2 = \{(c_1, \ldots, c_k) \mid \begin{array}{c} (c_1, \ldots, c_k) \in R_1 \text{ and} \\ (c_1, \ldots, c_k) \not\in R_2 \end{array} \} \]
$R_1 \times R_2 = \{(c_1, \ldots, c_k, d_1, \ldots, d_m) \mid (c_1, \ldots, c_k) \in R_1 \text{ and } (d_1, \ldots, d_m) \in R_2 \}.$
Let now $R$ be a relation of arity $k$ and $i_1, \ldots, i_m$ be numbers in \{1, \ldots, k\}.

$$
\pi_{i_1, \ldots, i_m}(R) = \{(c_{i_1}, \ldots, c_{i_m}) \mid (c_1, \ldots, c_k) \in R\}.
$$

We say that $\pi_{i_1, \ldots, i_m}(R)$ is obtained from $R$ by \textit{projection} (on arguments $i_1, \ldots, i_m$).
Assume **formulas on domains** with “variables” #1, #2, …. 

For example, #1 = #2.

\[
\sigma_F(R) = \{(c_1, \ldots, c_k) \mid (c_1, \ldots, c_k) \in R \text{ and } F \text{ holds on } (c_1, \ldots, c_k)\}.
\]
Overview

- Relational algebra, named perspective
- SQL
- Integrity constraints
- (Aggregates and grouping)
\[
\{ \{ a_1 = v_{11}, \ldots, a_n = v_{1n} \}, \]
\[
\ldots \ldots \ldots \ldots
\]
\[
\{ a_1 = v_{k1}, \ldots, a_n = v_{kn} \}\}
\]
Let $R_1$, $R_2$ be relations with the same attributes.

$$R_1 \cup R_2 = \{ t \mid t \in R_1 \text{ or } t \in R_2 \}$$
### Union, example

<table>
<thead>
<tr>
<th></th>
<th>R₁</th>
<th></th>
<th>R₂</th>
<th></th>
<th>R₁ ∪ R₂</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>A</strong></td>
<td><strong>B</strong></td>
<td><strong>A</strong></td>
<td><strong>B</strong></td>
<td><strong>A</strong></td>
<td><strong>B</strong></td>
</tr>
<tr>
<td>α</td>
<td>1</td>
<td>α</td>
<td>2</td>
<td>α</td>
<td>1</td>
</tr>
<tr>
<td>α</td>
<td>2</td>
<td>β</td>
<td>3</td>
<td>β</td>
<td>3</td>
</tr>
<tr>
<td>β</td>
<td>1</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Let $R$ be a relation whose set of attributes is $a_1, \ldots, a_n, c_1, \ldots, c_m$.

Let $b_1, \ldots, b_n$ be distinct attributes such that
\[
\{b_1, \ldots, b_n\} \cap \{c_1, \ldots, c_m\} = \emptyset
\]

Then
\[
\rho_{a_1 \rightarrow b_1, \ldots, a_n \rightarrow b_n}(R) = \\
\{\{b_1 = v_1, \ldots, b_n = v_n, c_1 = w_1, \ldots, c_m = w_m\} \mid \\
\{a_1 = v_1, \ldots, a_n = v_n, c_1 = w_1, \ldots, c_m = w_m\} \in R\}
\]
SQL is based on set and relational operations with certain modifications and enhancements

A typical SQL query has the form:

```
select a_1, \ldots, a_n
from R_1, \ldots, R_m
where P
```

This query is equivalent to relational algebra expression:

```
\pi_{a_1,\ldots,a_n}(\sigma_P(R_1 \times \ldots \times R_m))
```

The result of an SQL query is a relation.

Exceptions?
Integrity constraints

- Domain constraints.
- Key constraints.
- Foreign key constraints.

- More general, defined constraints.

- How to translate them?
Query language

Allow one to define:

- Relation and database schemas;
- Relations through our relations;
- Integrity constraints;
- Updates.