

## Models in First Order Logics



First-order logic. Syntax and semantics.  
Herbrand interpretations;  
Clauses and goals;  
Datalog.

First-order signature  $\Sigma$  consists of

- $con$  — the set of constants of  $\Sigma$ ;
- $fun$  — the set of function symbols of  $\Sigma$ ;
- $rel$  — the set of relation symbols of  $\Sigma$ .

Term of  $\Sigma$  with variables in  $X$ :

1. Constant  $c \in con$ ;
2. Variable  $v \in X$ ;
3. If  $f \in fun$  is a function symbol of arity  $n$  and  $t_1, \dots, t_n$  are terms, then  $f(t_1, \dots, t_n)$  is a term.

A term is ground if it has no variables

$\text{var}(t)$  — the set of variables of  $t$

Abstract notation:

- $a, b, c, d, e$  for constants;
- $x, y, z, u, v, w$  for variables;
- $f, g, h$  for function symbols;
- $p, q$  for relation symbols,

Example:  $f(x, g(y))$ .

Concrete notation:

Variable names start with upper-case letters.

Example:      likes(john, Anybody).

Atomic formulas, or atoms  $p(t_1, \dots, t_n)$ .

$(A_1 \wedge \dots \wedge A_n)$  and  $(A_1 \vee \dots \vee A_n)$

$(A \rightarrow B)$  and  $(A \leftrightarrow B)$

$\neg A$

$\forall v A$  and  $\exists v A$

Substitution  $\theta$  : is any mapping from the set  $V$  of variables to the set of terms such that there is only a finite number of variables  $v \in V$  with  $\theta(v) \neq v$ .

Domain  $dom(\theta)$ , range  $ran(\theta)$  and variable range  $vran(\theta)$ :

$$dom(\theta) = \{v \mid v \neq \theta(v)\},$$

$$ran(\theta) = \{ t \mid \exists v \in dom(\theta) (\theta(v) = t)\},$$

$$vran(\theta) = var(ran(\theta)).$$

Notation:  $\{ x_1 \mapsto t_1, \dots, x_n \mapsto t_n \}$

empty substitution  $\{\}$

Application of a substitution  $\theta$  to a term  $t$ :

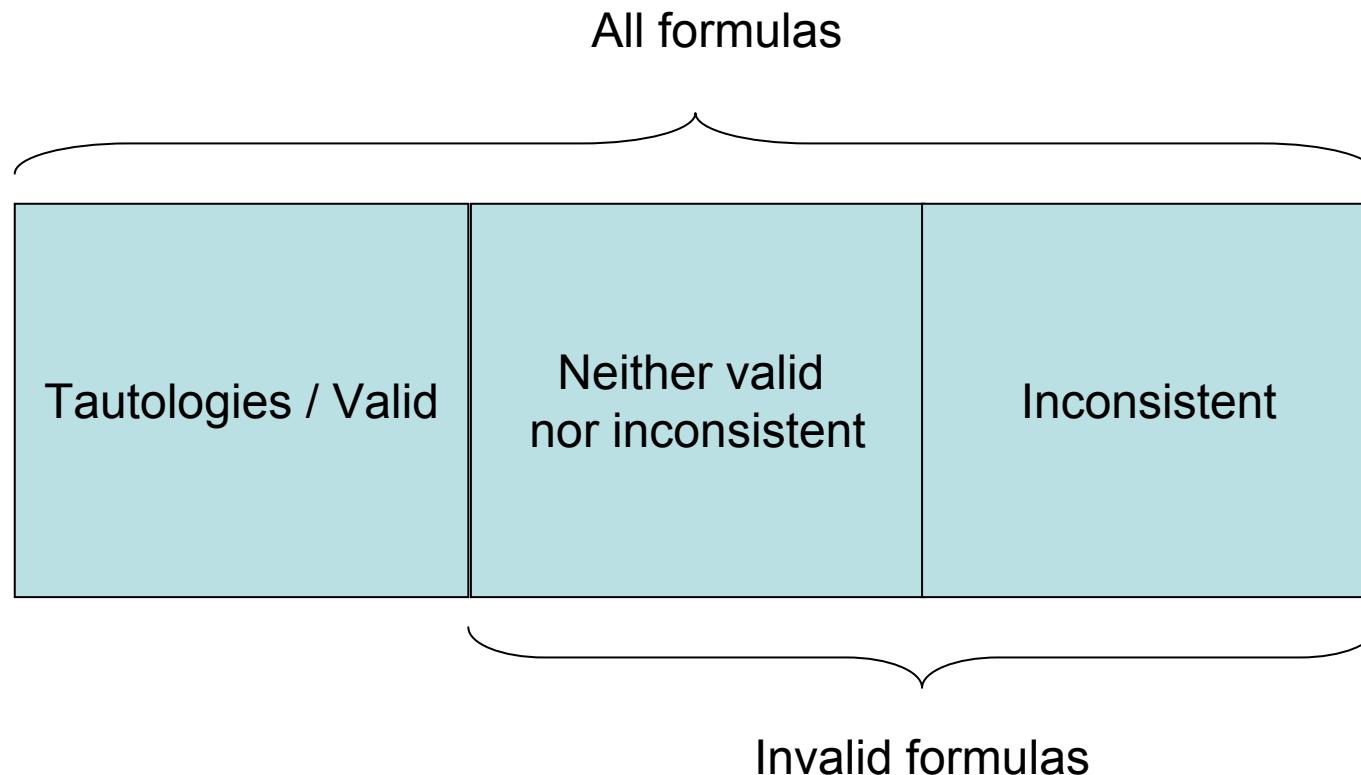
- $x\theta = \theta(x)$
- $c\theta = c$
- $f(t_1, \dots, t_n)\theta = f(t_1\theta, \dots, t_n\theta)$

A Herbrand interpretation of a signature  $\Sigma$  is a set of ground atoms of this signature.

1. If  $A$  is atomic, then  $I \models A$  if  $A \in I$
2.  $I \models B_1 \wedge \dots \wedge B_n$  if  $I \models B_i$  for all  $i$
3.  $I \models B_1 \vee \dots \vee B_n$  if  $I \models B_i$  for some  $i$
4.  $I \models B_1 \rightarrow B_2$  if either  $I \models B_2$  or  $I \not\models B_1$
5.  $I \models \neg B$  if  $I \not\models B$
6.  $I \models \forall x B$  if  $I \models B\{x \mapsto t\}$  for all ground terms  $t$  of the signature  $\Sigma$
7.  $I \models \exists x B$  if  $I \models B\{x \mapsto t\}$  for some ground term  $t$  of the signature  $\Sigma$

A formula  $F$  is a tautology (is valid), if  $I \models F$  for every (Herbrand) interpretation  $I$

A formula  $F$  is inconsistent, if  $I \not\models F$  for every (Herbrand) interpretation  $I$ .



A (set of) formula(s)  $F$  logically implies  $G$  (we write  $F \models G$ ), iff every (Herbrand) interpretation  $I$  that fulfills  $F$  ( $I \models F$ ) also fulfills  $G$  ( $I \models G$ ).

$\wedge F \models G$  is true iff for every (Herbrand) interpretation  $I$ :  
 $I \models (\wedge F \rightarrow G)$

Literal is either an atom or the negation  $\neg A$  of an atom  $A$ .

Positive literal: atom

Negative literal: negation of an atom

Complementary literals:  $A$  and  $\neg A$

Notation: L

Clause: (or normal clause) formula  $L_1 \wedge \dots \wedge L_n \rightarrow A$ ,  
where

$n \geq 0$ , each  $L_i$  is a literal and  $A$  is an atom.

Notation:  $A :- L_1 \wedge \dots \wedge L_n$  or  $A :- L_1, \dots, L_n$

Head: the atom  $A$ .

Body: The conjunction  $L_1 \wedge \dots \wedge L_n$

Definite clause: all  $L_i$  are positive

Fact: clause with empty body

Clause	Class
lives(Person, sweden) :- sells(Person, wine, Shop), not open(Shop,saturday)	normal
spy(Person) :- russian(Person)	definite
spy(bond)	fact

- Goal (also normal goal) is any conjunction of literals  
 $L_1 \wedge \dots \wedge L_n$
- Definite goal: all  $L_i$  are positive
- Empty goal  $\square$ : when  $n = 0$

## Excercise:

- Syntax & Semantics for Monadic Fuzzy Logics
- Monadic: only unary predicates
- Fuzzy:
  - ◆ Truth values in  $[0,1]$
  - ◆ operators for truth values
    - T-norm:  $\top_{\min}(a,b)=\min\{a,b\}$
    - T-conorm:  $\perp_{\min}(a,b)=\max\{a,b\}$
    - Also:  $\top_{\text{prod}}(a,b)=a \cdot b$ ,  $\perp_{\text{sum}}(a,b)=a+b-a \cdot b, \dots$
    - Duality for  $N(x)=1-x$